

**RYERSON UNIVERSITY**  
**DEPARTMENT OF COMPUTER SCIENCE**

**CPS 616**  
**FINAL EXAM**  
**WINTER 2015**

NAME: \_\_\_\_\_

STUDENT ID: \_\_\_\_\_

**INSTRUCTIONS**

- This exam is 3 hours long.
- This exam is out of 70 and is worth 35% of the course mark.
- This is a closed book exam. However, one double-sided letter-sized crib sheet is allowed.
- This exam is single-sided and has 6 pages including this front page.
- The first part of this exam is multiple choice. Please enter your answers to this part on the bubble sheet provided.
- The second part of this exam is short answer. Please answer all questions of this part directly on this exam.

**PART 1 - MULTIPLE CHOICE**

**Instructions**

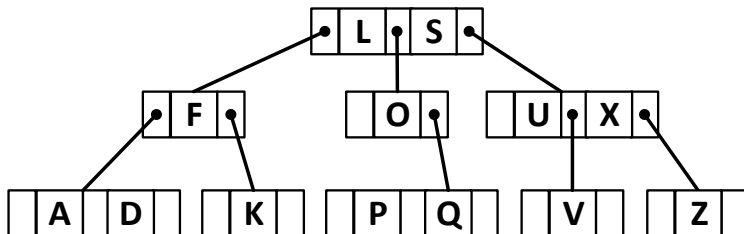
Please enter your answers on the bubble sheet with your name.

Use pencil only.

You may guess. Marks will not be deducted for false answers..

**Questions (15 marks)**

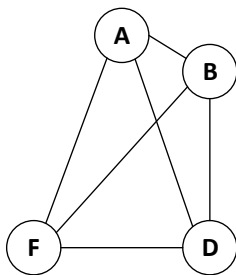
1. (2 marks) Is this graph:



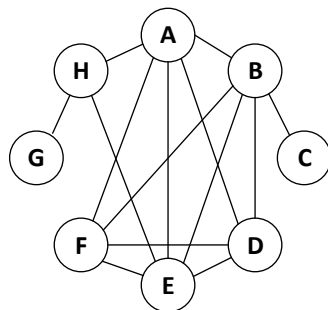
- A. An AVL Tree
- B. A 2-3 Tree
- C. A Red-Black Tree
- D. None of the above

2. (2 marks) Given the following graphs, which statement is correct:

Graph G



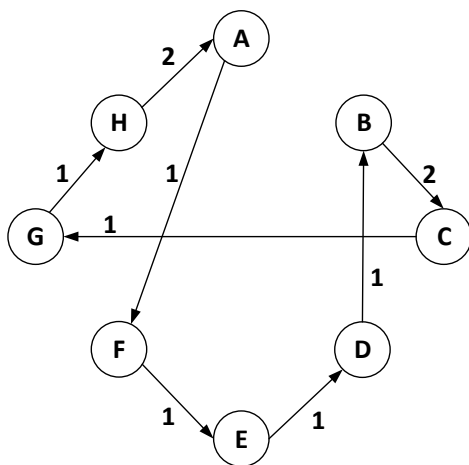
Graph H



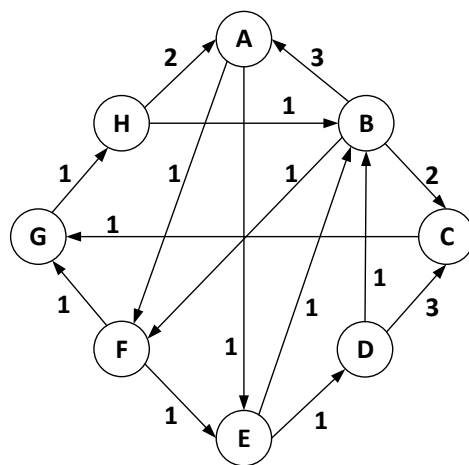
- A. G is the complement of H
- B. G is a non-maximal clique of H
- C. G is a maximal clique of H
- D. G is a spanning tree of H
- E. None of the above

3. (3 marks) Given the following graphs, which statement is correct:

Graph F



Graph G



- A. F is an Eulerian circuit of G which **is not** a solution to the Travelling Salesman problem for G
- B. F is a Hamiltonian circuit of G which **is not** a solution to the Travelling Salesman problem for G
- C. F is an Eulerian circuit of G which **is** a solution to the Travelling Salesman problem for G
- D. F is a Hamiltonian circuit of G which **is** a solution to the Travelling Salesman problem for G
- E. None of the above

The next 2 questions refer to the following linear programming problem whose feasible region and solution is graphed below. **The two questions are independent of each other.**

Linear Programming problem

Maximise

$$x+y$$

For

$$0 \leq x$$

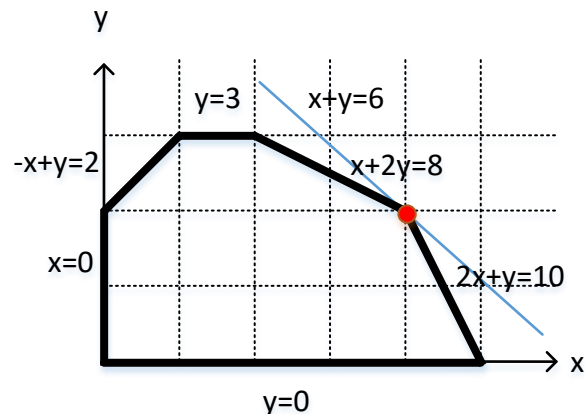
$$0 \leq y \leq 3$$

$$-x+y \leq 2$$

$$x+2y \leq 8$$

$$2x+y \leq 10$$

Feasible region and solution



4. (2 marks) What happens to this problem if you change the constraint  $y \leq 3$  to be  $y \geq 3$ ?
  - A. The problem becomes unfeasible and has no solution
  - B. The problem becomes unbounded and has no solution
  - C. The problem remains feasible and has a single optimal solution
  - D. The problem remains feasible but has more than one optimal solution
5. (2 marks) What happens to this problem if you change the constraint  $y \leq 3$  to be  $y \geq 5$ ?
  - A. The problem becomes unfeasible and has no solution
  - B. The problem becomes unbounded and has no solution
  - C. The problem remains feasible and has a single optimal solution
  - D. The problem remains feasible but has more than one optimal solution
6. (2 marks) Which of the following statements is **not true** about the following problem P:  
find the maximum element of an  $n \times n$  array
  - A.  $\Omega(n)$  is a lower bound for P
  - B.  $\Omega(n)$  is a tight lower bound for P
  - C.  $\Omega(n^2)$  is a lower bound for P
  - D.  $\Omega(n^2)$  is a tight lower bound for P
7. (2 marks) It can be proven that for two positive integers  $m$  and  $n$  such that  $m \leq n$  the worst case cost of the Euclidian algorithm to calculate  $\gcd(m,n)$  is  $O(\log_{10}m)$ .

Also, as you know,  $\text{lcm}(m,n) = m.n / \gcd(m,n)$

What can we therefore conclude about the cost  $C$  of calculating  $\text{lcm}(m,n)$  using the Euclidian algorithm?

- A.  $C \in O(\log_{10}m)$
- B.  $C \in O(1/\log_{10}m)$
- C.  $C \in O(n/\log_{10}m)$
- D.  $C \in O(n.m/\log_{10}m)$

**PART 2 - SHORT ANSWERS - PLEASE WRITE YOUR ANSWERS DIRECTLY IN THIS EXAM**

**8. (20 marks) - Huffman Code**

- (16 marks) Construct a Huffman tree and corresponding Huffman code for the following data. Your internal nodes should contain total frequencies, and your branches should be labelled with 0s and 1s

Symbol	' '	a	b	d	e	g	o	y
Frequency	0.20	0.15	0.11	0.13	0.17	0.10	0.10	0.04
Codeword								

When there are ties for the selection of the subtrees that will form a tree, they are broken in accordance with the rules below in the following order:

- 1) The shallowest subtree(s) is(are) selected
- 2) The subtree(s) containing the earliest element in alphabetical order is(are) selected

When building a tree from 2 subtrees, the left and right subtrees are placed in accordance with the rules below in the following order:

- 1) The lighter subtree is on the left
- 2) The shallowest subtree is on the left
- 3) The tree containing the earliest element in alphabetical order is on the left

(2 marks) Encode the string "**doggy bag**" (2 words with a blank between them) using the Huffman code you designed in a)

(2 marks) Decode the word 10110100101000111010111 using the Huffman code you designed in a)

**9. (20 marks) Non-Deterministic Turing Machines (NTMs)**

Here is a NTM with 3 tapes called T1,T2,T3.

The symbol "-" is the blank symbol.

The actions for each tape have the format: <new symbol> , <move>

where the possible moves are: R=move right, L=move left, S=stay

State	Current Symbol			Action			New State	Explanation
	T1	T2	T3	T1	T2	T3		
q <sub>1</sub>	1	-	-	1,R	1,R	-,S	q <sub>2</sub>	
	1	-	-	1,R	-,S	1,R	q <sub>3</sub>	
q <sub>2</sub>	0	-	-	0,R	0,R	-,S	q <sub>2</sub>	
	1	-	-	1,S	-,S	-,S	q <sub>1</sub>	
	-	-	-	-,S	-,L	-,L	q <sub>4</sub>	
q <sub>3</sub>	0	-	-	0,R	-,S	0,R	q <sub>3</sub>	
	1	-	-	1,S	-,S	-,S	q <sub>1</sub>	
	-	-	-	-,S	-,L	-,L	q <sub>4</sub>	
q <sub>4</sub>	-	0	0	-,S	0,L	0,L	q <sub>4</sub>	
	-	0	1	-,S	0,S	1,L	q <sub>4</sub>	
	-	1	0	-,S	1,L	0,S	q <sub>4</sub>	
	-	1	1	-,S	1,L	1,L	q <sub>4</sub>	
	-	-	-	-,S	-,S	-,S	q <sub>5</sub>	
q <sub>5</sub>								Accept

- (1 mark) Circle all the non-deterministic states in the first column of the table

(6 marks) Explain in English in the last column of the table what the states q<sub>1</sub> to q<sub>4</sub> do.

(3 marks) Why does state q<sub>5</sub> accept the input on T1? In other words, what does this NTM do?

(2 marks) If this NTM accepts an input string on T1 which has  $n$  1's, how many times will it have read a 1 before it accepted the string? The answer should be a function of  $n$ . In this question "reading a 1" means that one of the tape heads was over a 1 and moved left or right away from it.

(3 marks) Why is the previous question asking you to count how many times a 1 is read and not to count how many times a 0 is read?

(3 marks) Is the problem that this NTM solves in NP class? Explain your answer.

(2 marks) Is the problem that this NTM solves in P class? Explain your answer.

**10. (15 marks) Dynamic Programming**

Here is pseudocode for a recursive program

```
// This function calculates x to the power of n
// You can assume that x is positive and n is a non-negative integer
exponent(x, n)
    if n=0 return 1
    if n=1 return x
    return exponent(x, n div 2) * exponent(x, n - (n div 2))
```

- (2 marks) Assuming that the multiplication in the last line is the basic operation in this algorithm, what is the exact cost of this algorithm as a function of  $n$ ?

(8 marks) Modify this pseudocode so that it solves exactly the same problem by recursing exactly the same way, but using a dynamic programming approach, i.e:

- "exponent(x, n div 2) \* exponent(x, n - (n div 2))" will still be part of your code
- solutions to subproblems are stored the first time they are calculated and never calculated again.

You can use the following global array:

```
global exp[0,..,n]    // to store intermediate results
                    // This array is initialized with 0s
```

- (2 marks) What is the best case cost of your new algorithm as a function of  $n$ ?
- (3 marks) Give an upper bound for the worst case cost of your new algorithm as a function of  $n$ . (Hint: look at  $n=127$ )  
**Do not spend too much time on this question, it is only worth 3 marks!**